

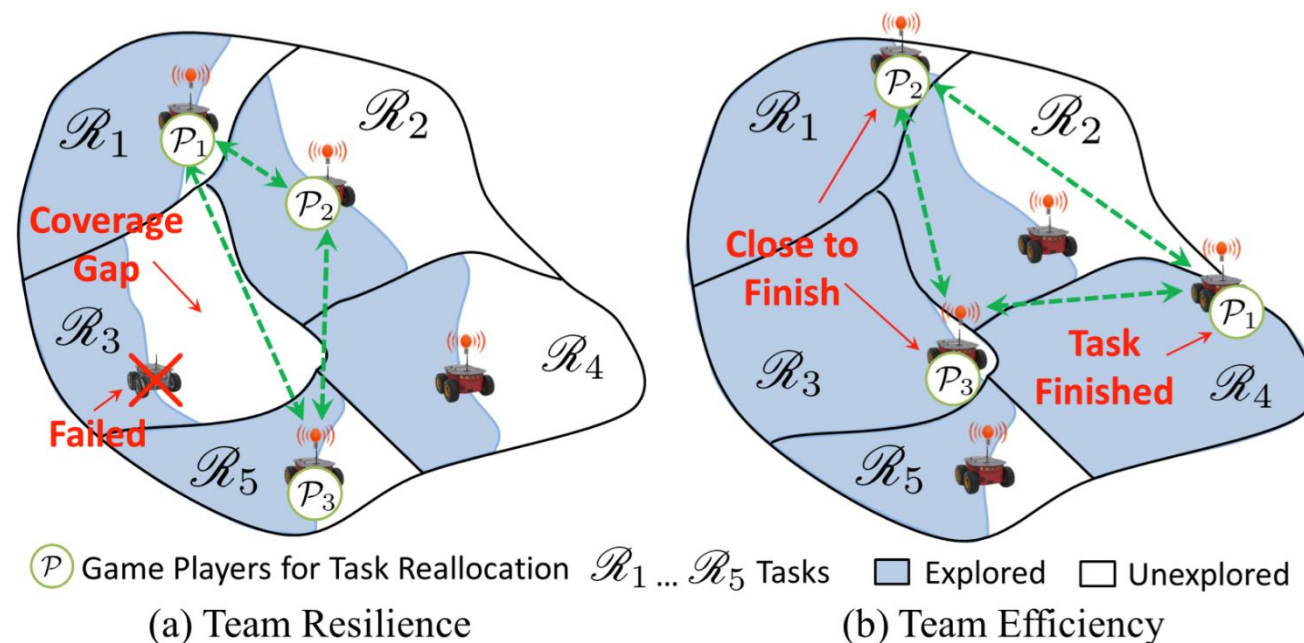
CARE: Cooperative Autonomy for Resilience and Efficiency of Robot Teams for Complete Coverage in Unknown Environment

Authors: Junnan Song and Shalabh Gupta

❖ **Objective:** Develop a multi-robot resilient and efficient algorithm for complete coverage in unknown environment.

❖ **Challenges:**

- Scalability: distributed vs. centralized
- Dynamically changing conditions
- Resilience: complete coverage under failures
- Efficiency: prevent robot idling
- Optimization factors during task reallocations:
 - Task worth (e.g., unfound targets)
 - Probabilities of success of the available robots in finishing the contested tasks
- Connection between local and global objectives: the local optimization must not only benefit the involved robots but also the whole team
- Real-time execution



Requirements for resilience and efficiency

The Autonomous Vehicle and the Search Area

❖ The Robot Team $V = \{v_\ell, \ell = 1, \dots, N\}$

1. Localization system, range detector and tasking sensor.
2. Wireless Communication Device
 - Allows periodic information exchange between all pairs of robots. The communication is assumed to be perfect.

❖ Battery Reliability

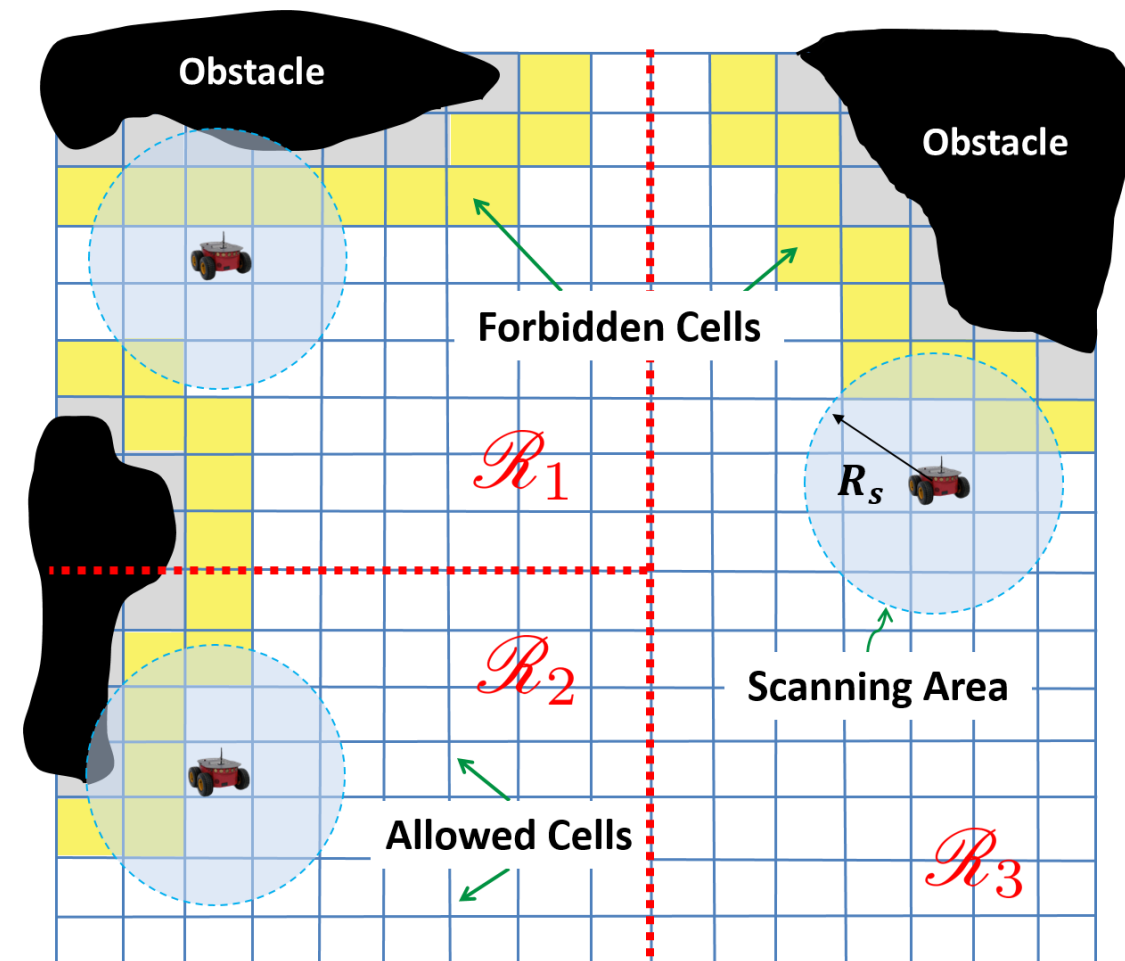
Each robot $v_\ell \in V$ is assumed to carry a battery, whose reliability is computed as^{[1][2]}

$$R_{v_\ell}(t) = \frac{1}{1 + e^{\rho_0(t - \rho_1)}}$$

where ρ_0 and ρ_1 are determined based on specific batteries.

❖ Initial Task Allocation

The tiling is grouped into M disjoint regions $\{\mathcal{R}_r, r = 1, \dots, M\}$, s.t. $\mathcal{R} = \bigcup_{r=1}^M \mathcal{R}_r$. Each robot can work on one task at a time, but one task can be assigned to multiple robots.



An example of the area \mathcal{R} with $M = 3$

[1] A. Islam, A. Alim, C. Hyder, and K. Zubaer, "Digging the inner reliability of wireless networked systems," in *2015 International Conference on Networking Systems and Security*, pp. 1–10, IEEE, 2015.

[2] M. Jongerden and B. Haverkort, "Which battery model to use?," *IET software*, vol. 3, no. 6, pp. 445–457, 2009.

❖ Complete Coverage

Let $\epsilon_\ell(k) \in T$ be the ϵ -cell that is visited and explored by robot v_ℓ at time k . Then the team V is said to achieve complete coverage, if $\exists K \in \mathbb{N}$, s.t. the sequences $\{\epsilon_\ell(k), k = 1, \dots, K\}, \forall \ell = 1, \dots, N$, satisfy

$$\mathcal{R}(T^a) \subseteq \bigcup_{\ell=1}^N \bigcup_{k=1}^K \epsilon_\ell(k)$$

❖ Performance Metrics

- Coverage Ratio (**CR**):

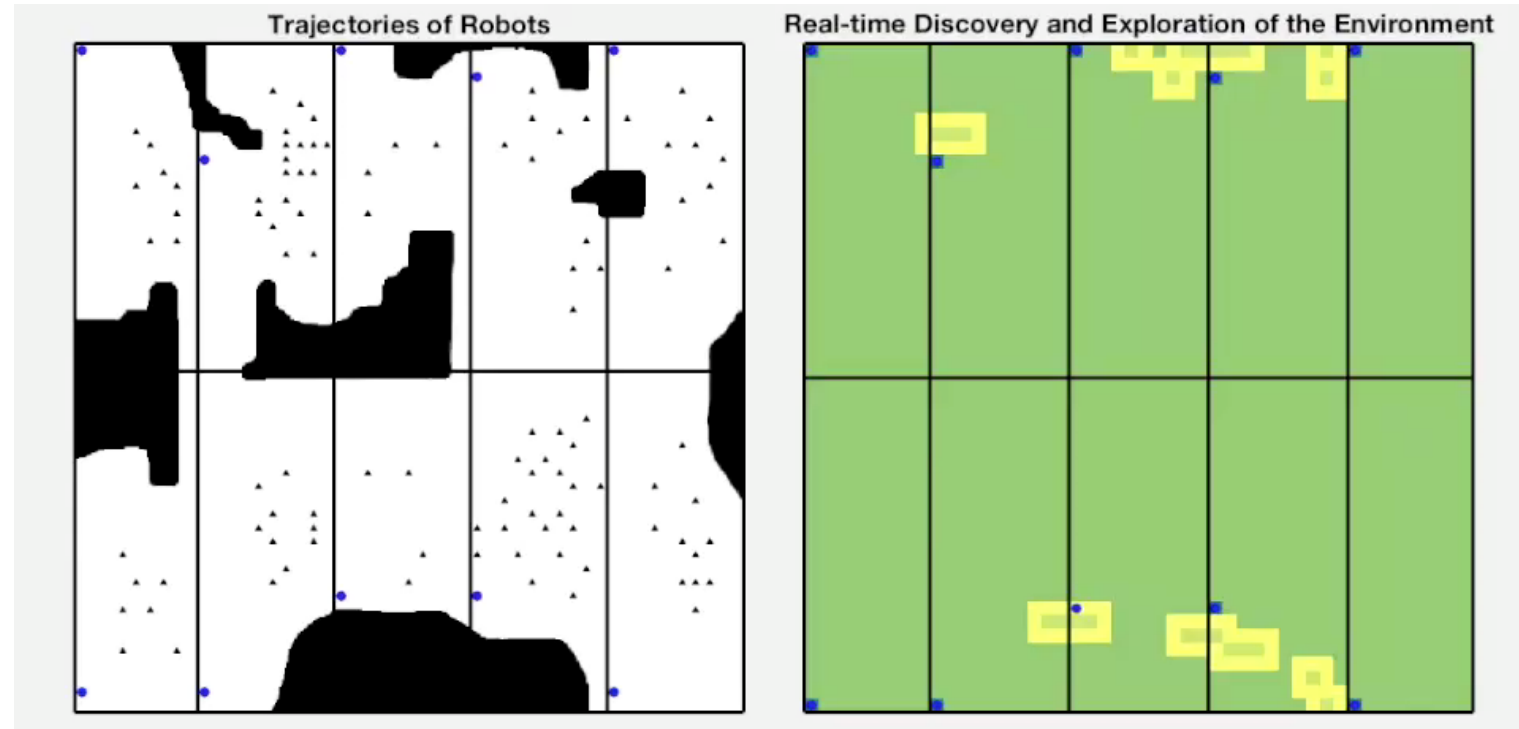
$$CR = \frac{(\bigcup_{\ell=1}^N \bigcup_{k=1}^K \epsilon_\ell(k)) \cap \mathcal{R}(T^a)}{\mathcal{R}(T^a)} \in [0,1]$$

- Coverage Time (**CT**): measured by the last finishing robot
- Remaining Reliability (**RR**): the average remaining reliability of all live robots by the end of the operation
- Number of Targets Found (**NoTF**): total number of targets discovered by the whole team
- Time of Target Discovery (**ToTD**): time for the team to discover a certain percentage of all targets

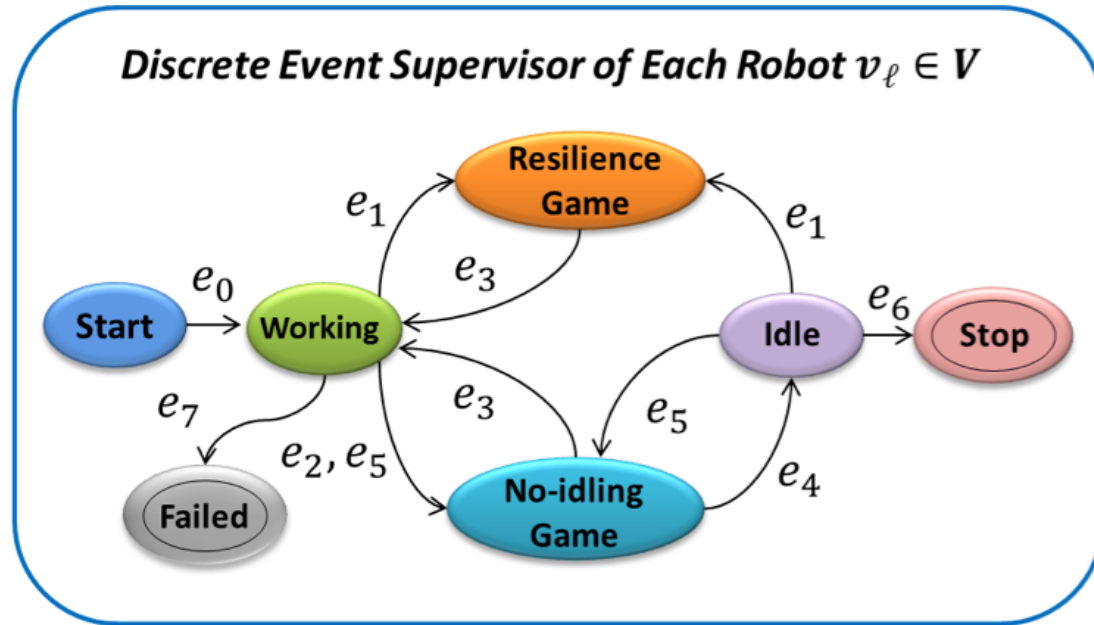
Objective: To achieve $CR = 1$ (even under a few robot failures), while minimizing CT and $ToTD$, and maximizing RR and $NoTF$.

❖ Features:

- Distributed multi-robot control
- Complete coverage under failures
- Proactive event-driven task reallocation upon robot failures or robot idling
- Task reallocation considers various optimization factors, including task worth, robot remaining energy, and relative locations.
- Local task reallocation decision is aligned with the global objective of the whole team



Distributed Discrete Event Supervisory Control Structure



Events	Description	Generation Condition
e_0	Start exploration	Robot v_ℓ is turned on
e_1	Neighbor failed	Confirmation of neighbor failure
e_2	Own task completed	$n_U(r_c(v_\ell)) = 0$
e_3	Task assigned	Task assigned by the Optimizer
e_4	No task assigned	No task assigned by the Optimizer
e_5	Neighbor task completed	Some neighbor completes its task and $t_c(r_c(v_\ell)) \leq \eta$
e_6	All tasks completed	$\sum n_U(r) = 0$
e_7	Robot failed	Robot v_ℓ is diagnosed as failed

- **Failure Detection:** use a standard mechanism of heartbeat signals^[1], where each robot v_ℓ periodically broadcasts heartbeat signals indicating its healthiness, and also listening from others.
- A robot is detected as **failed** if its heartbeat signal cannot be received constantly for a certain period of time.

- $r_c: V \rightarrow \{1, \dots, M\}$ is the allocation function that indicates the current task allocations of robots
- $t_c: \{1, \dots, M\} \rightarrow [0, \infty)$ is the remaining estimated time to finish a given task by its assigned robots
- $n_U: \{1, \dots, M\} \rightarrow \mathbb{N}$ is the number of unexplored cells in a task
- $\eta \in \mathbb{R}^+$: the threshold to define a task as *close-to-finish*.

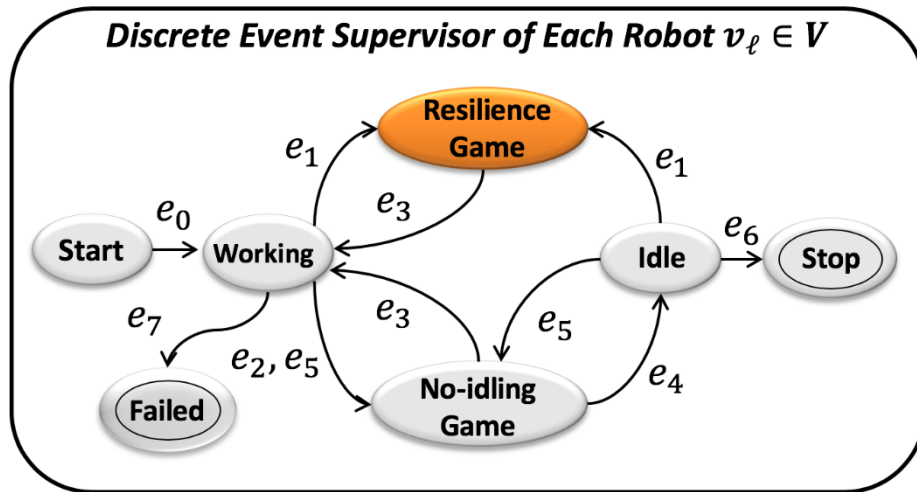
[1] W. Chen, S. Toueg, and M. Aguilera, "On the quality of service of failure detectors," *IEEE Transactions on Computers*, vol. 51, no. 5, pp. 561–580, 2002.

When is it triggered?

Upon detection of a neighbor failure (i.e., e_1)

What does it do?

Re-optimize local task allocations to fill coverage gap.

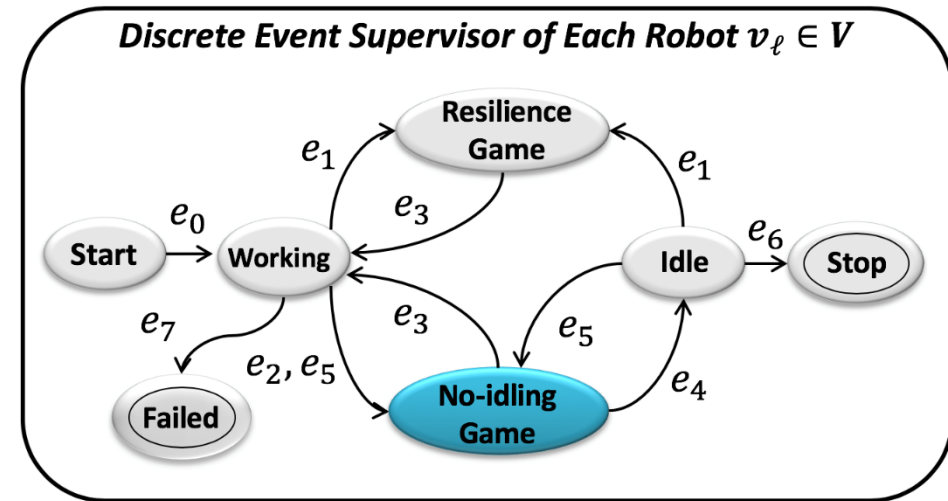


When is it triggered?

Upon completion of own task (i.e., e_2), or a neighbor's task (i.e., e_5)

What does it do?

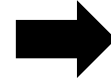
Compute for new tasks (if available).



Task re-allocation is required!

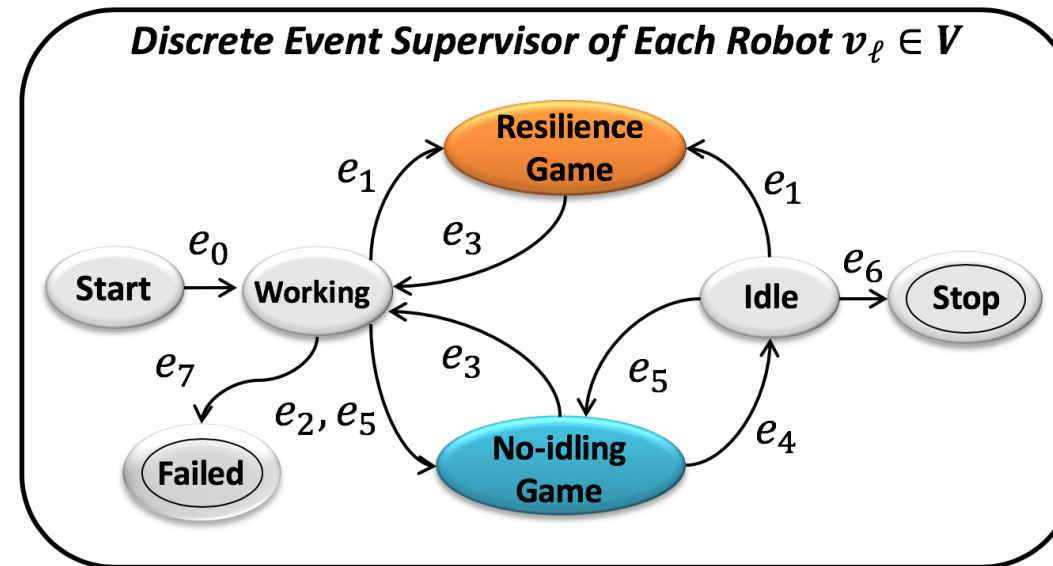
❖ Requirements for Task Reallocation:

- Distributed framework for scalability
- Optimization over both task worth and robot reliabilities
- Local reallocation decision must also benefit the whole team
- Complete coverage under failures



❖ Potential Games

- Fits the distributed framework
- Utility of each player is perfectly aligned with the global potential function ϕ for the whole team.
- Existence of solutions: at least one pure Nash Equilibrium exists, which is a maximizer to ϕ
- Max-Logit can quickly converge to the (sub-)optimal equilibrium



Game Components:

- Players $\mathcal{P} = \{\mathcal{P}_i \in V, i = 1, \dots, |\mathcal{P}|\}$
 - The available robots to be reallocated.
- Action set \mathcal{A}_i for each player \mathcal{P}_i
 - Each action $a_i \in \mathcal{A}_i$ refers to the index of an available task. Also, it has $\tilde{\mathcal{A}} = \mathcal{A}_i = \mathcal{A}_j, \forall i \neq j$.

- The *utility function* for each player \mathcal{P}_i , defined as:

$$U_i: \mathcal{A}_{\mathcal{P}} \rightarrow \mathbb{R}$$

where $\mathcal{A}_{\mathcal{P}} = \mathcal{A}_1 \times \mathcal{A}_2 \times \dots \times \mathcal{A}_{|\mathcal{P}|}$ denotes the set of joint actions for all players, and the joint action $a_{\mathcal{P}} = (a_i, a_{-i}) \in \mathcal{A}_{\mathcal{P}}$ indicates a task reallocation for the team.

- Note: players other than \mathcal{P}_i have joint action $a_{-i} \in \mathcal{A}_{-i}$

Nash Equilibrium: A joint action $a_{\mathcal{P}}^* = (a_i^*, a_{-i}^*) \in \mathcal{A}_{\mathcal{P}}$ is a pure Nash Equilibrium if:

$$U_i(a_i^*, a_{-i}^*) = \max_{a_i \in \tilde{\mathcal{A}}} U_i(a_i, a_{-i}^*), \forall \mathcal{P}_i \in \mathcal{P}$$

Potential Game: a potential game G in strategic form is a potential game if and only if a potential function $\phi: \mathcal{A}_{\mathcal{P}} \rightarrow \mathbb{R}$ exists, s.t., $\forall \mathcal{P}_i \in \mathcal{P}$

$$U_i(a'_i, a_{-i}) - U_i(a''_i, a_{-i}) = \phi(a'_i, a_{-i}) - \phi(a''_i, a_{-i})$$

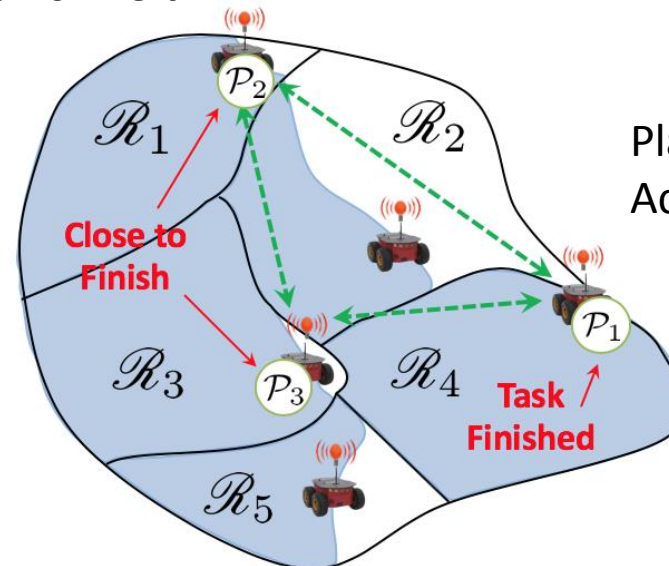
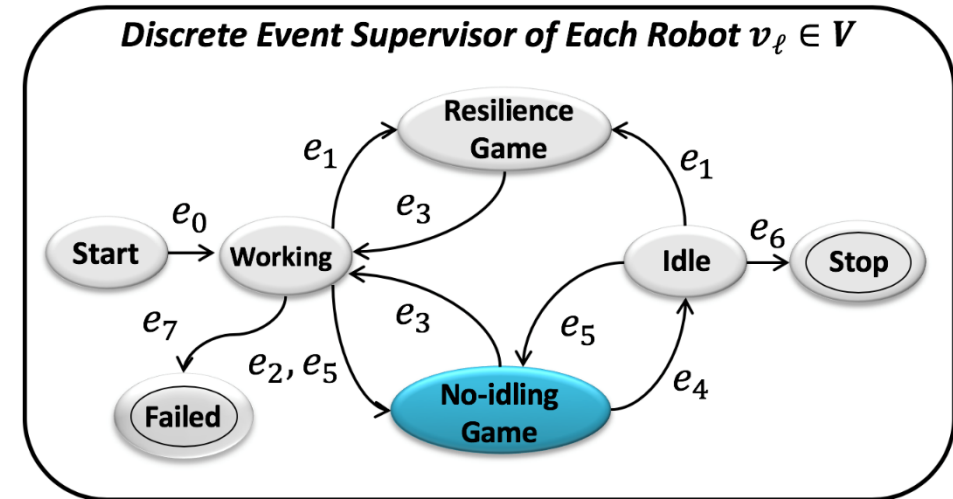
$\forall a'_i, a''_i \in \mathcal{A}_i$ and $a_{-i} \in \mathcal{A}_{-i}$.

The CARE Algorithm

Specifics of the No-idling Games

❖ No-idling Games

- **Condition:** if some robot $v_{id} \in V$ completes its current task and becomes idle.
- **Game Specifics:** It calls the κ_1 nearest neighbors $N_{\kappa_1}^{v_{id}}$, that are close to finish their tasks to participate.
 - Players: $\mathcal{P} = \{v_{id}\} \cup \{v_\ell \in N_{\kappa_1}^{v_{id}} : t_c(r_c(v_\ell)) \leq \eta\}$
 - Actions: $\mathcal{A}_i = \{r \in \{1, \dots, M\} : t_c(r) \geq \gamma \in \mathbb{R}^+\}$, i.e., the incomplete tasks with sufficient work left.



Players: $\mathcal{P}_1, \mathcal{P}_2, \mathcal{P}_3$
 Actions: $\mathcal{R}_2, \mathcal{R}_5$

\mathcal{R} Tasks \mathcal{P} Players Explored Unexplored

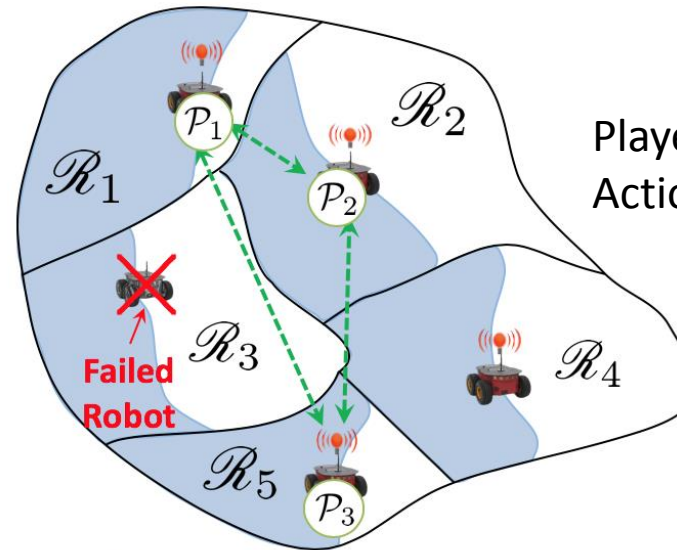
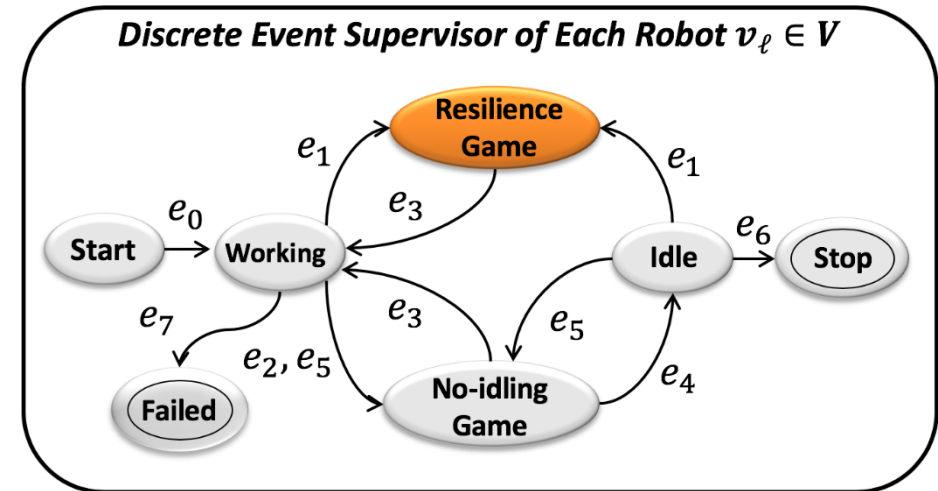
Figure. An example of the no-idling game when $\kappa_1 = 3$

The CARE Algorithm

Specifics of the Resilience Games

❖ Resilience Games

- **Condition:** if some robot $v_f \in V$ fails.
- **Game Specifics:** the κ_2 nearest neighbors of v_f , $N_{\kappa_2}^{v_f}$, are involved.
 - **Players:** $\mathcal{P} = N_{\kappa_2}^{v_f}$
 - **Actions:** $\mathcal{A}_i = \{r_c(v_f)\} \cup \{r_c(v_\ell), v_\ell \in N_{\kappa_2}^{v_f} : t_c(r_c(v_\ell)) \geq \eta\}$,
i.e., the current tasks of all players and the failed robot.



Players: $\mathcal{P}_1, \mathcal{P}_2, \mathcal{P}_3$
 Actions: $\mathcal{R}_1, \mathcal{R}_2, \mathcal{R}_3, \mathcal{R}_5$

\mathcal{R} Tasks \mathcal{P} Players Explored Unexplored

Figure. An example of the resilience game when $\kappa_2 = 3$

The CARE Algorithm

Potential Function and Utility Function

❖ Potential Function $\phi(a_{\mathcal{P}})$

$$\phi(a_{\mathcal{P}}) = \sum_{r \in \mathcal{A}} w_r \left(1 - \prod_{\mathcal{P}_i \in \{\mathcal{P}\}_r} [1 - p_r(\mathcal{P}_i)] \right)$$

Joint probability to successfully finish task r

where:

- w_r : estimated worth (i.e., undiscovered number of targets) of task $r \in \mathcal{A}_i$
- $\{\mathcal{P}\}_r = \{\mathcal{P}_i \in \mathcal{P} : a_i = r\}$: set of players that choose task r in the joint action $a_{\mathcal{P}}$
- $p_r(\mathcal{P}_i)$: the success probability for player \mathcal{P}_i to finish task r

Physical Meaning of $\phi(a_{\mathcal{P}})$: the total expected worth for the joint action $a_{\mathcal{P}} \in \mathcal{A}_{\mathcal{P}}$

❖ Utility Function

The utility function is obtained using *Marginal Contribution*:

$$\begin{aligned} \mathcal{U}_i(a_i, a_{-i}) &= \phi(a_i, a_{-i}) - \phi(\emptyset, a_{-i}) \\ &= w_{a_i} \cdot p_{a_i}(\mathcal{P}_i) \cdot \prod_{\mathcal{P}_j \in \{\mathcal{P}\}_{a_i} \setminus \mathcal{P}_i} [1 - p_{a_i}(\mathcal{P}_j)] \end{aligned}$$

Where \emptyset represents the null action.

Joint failure probability for task r by other players

Proposition 1: The game with potential function ϕ and utility function \mathcal{U}_i constitute a potential game^[1].

[1] J. Song and S. Gupta, "CARE: Cooperative Autonomy for Resilience and Efficiency of Robot Teams for Complete Coverage of Unknown Environment under Robot Failures," *Autonomous Robots*, Under review, 2018.

Success Probability $p_r(\mathcal{P}_i)$: the probability for player \mathcal{P}_i to successfully finish a contested task r

❖ Compute $p_r(\mathcal{P}_i)$

It is evaluated based on robot reliability

$$p_r(\mathcal{P}_i) = R_{\mathcal{P}_i}(\tilde{t})$$

Where:

- Battery reliability: $R_{\mathcal{P}_i}(t) = \frac{1}{1+e^{\rho_0(t-\rho_1)}}$
- Expected total time of tasking and traveling:

$$\tilde{t} = t_k + t_{tr} + t_r$$

- t_k : the total tasking time of \mathcal{P}_i since the beginning until game is initiated.
- $t_{tr} = \frac{Dist(\mathcal{P}_i, r)}{u}$: the traveling time to task r , and $u \in \mathbb{R}^+$ is the traveling speed.
- $t_r = \frac{N_U(r)}{\omega}$: the estimated time to finish task r , and $\omega \in \mathbb{R}^+$ is the speed of tasking a cell.
- Additionally, if a player has small portion of work left in its current task (i.e., $t_c(r_c(\mathcal{P}_i)) \leq \eta$), then t_c is also included into \tilde{t} .

Task Worth w_r : the expected number of undiscovered targets in task r that are available to the players.

❖ Total Worth of Task r

- Let x_r be a random variable that denotes the total number of targets in task r . It is assumed to follow the *Poisson distribution* with a *known* mean λ_r :

$$P_r(x_r = x) = e^{-\lambda_r} \cdot \frac{\lambda_r^x}{x!}, x = 0, 1, 2, \dots$$

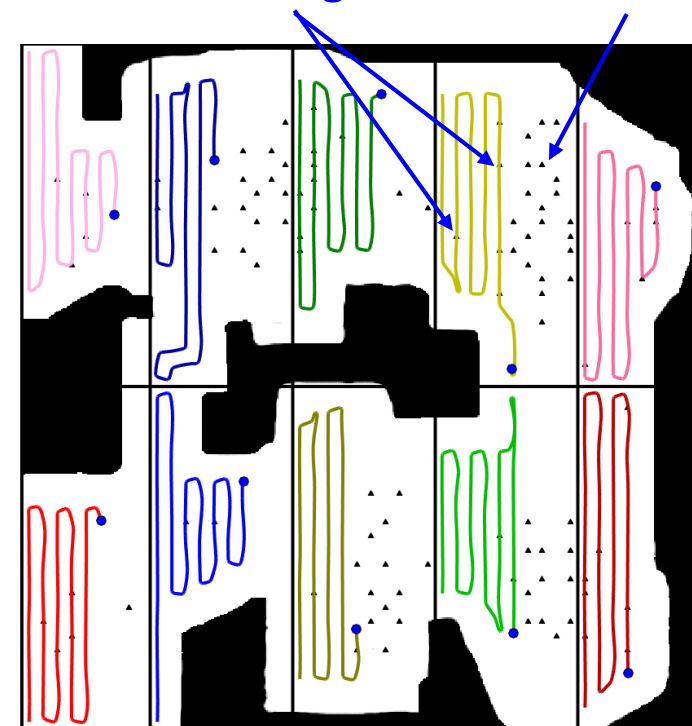
- If ξ have already been discovered, then the estimated remaining number of undiscovered targets are:

$$\bar{w}_r = \sum_{x=\xi+1}^{\infty} (x - \xi) \cdot e^{-\lambda_r} \cdot \frac{\lambda_r^x}{x!}$$

Using $\sum_{x=0}^{\infty} x \cdot e^{-\lambda_r} \cdot \frac{\lambda_r^x}{x!} = \lambda_r$, and $\sum_{x=0}^{\infty} e^{-\lambda_r} \cdot \frac{\lambda_r^x}{x!} = 1$, then:

$$\bar{w}_r = (\lambda_r - \xi) + e^{-\lambda_r} \cdot \sum_{x=0}^{\xi} (\xi - x) \cdot \frac{\lambda_r^x}{x!}$$

discovered targets undiscovered targets



Robot exploration with target discovery

The CARE Algorithm

Task Worth w_r

But... there can be non-player robots currently working in task r , and they are also discovering targets but not participating the game.

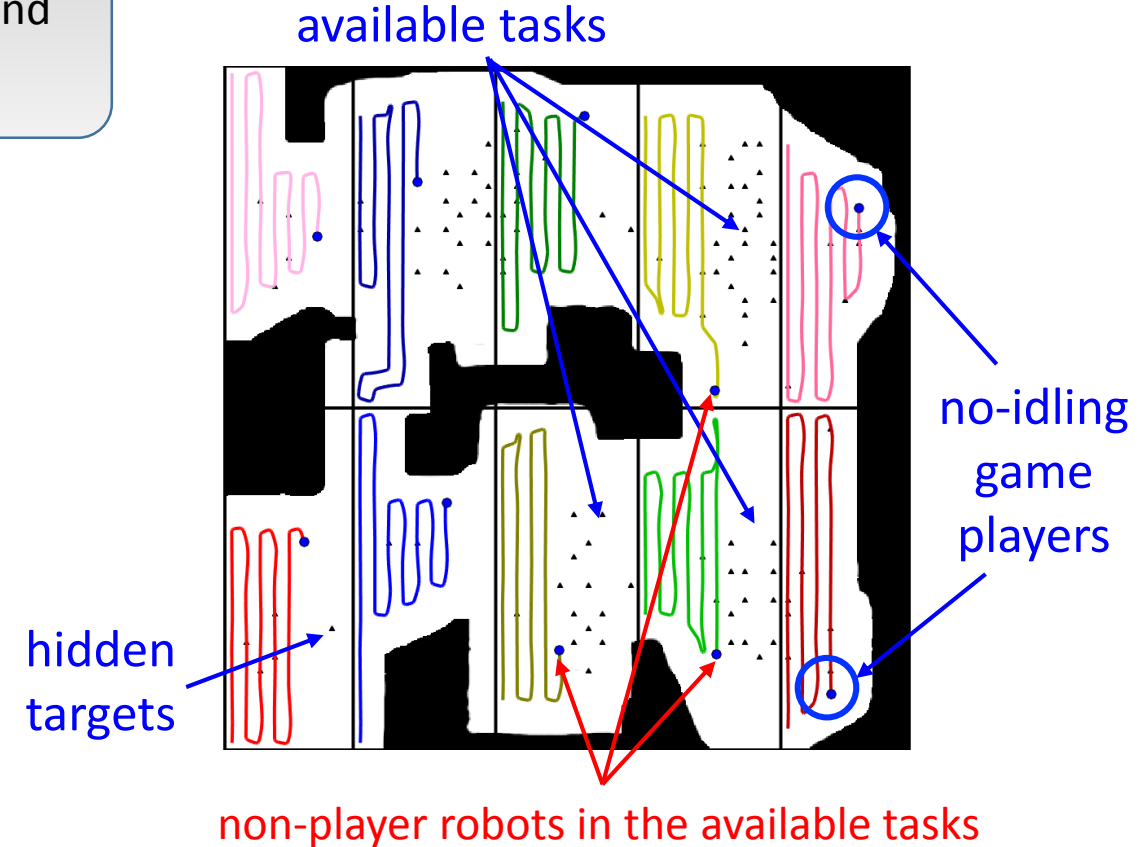
❖ Task Worth w_r for Players

- Let $\bar{\mathcal{P}} = V \setminus \mathcal{P}$ denote the set of non-players.
- Let $\{\bar{\mathcal{P}}\}_r$ denote the set of non-players currently in task r .
- These non-players have joint success probability for task r :

$$q(r) = 1 - \prod_{v_\ell \in \{\bar{\mathcal{P}}\}_r} [1 - p_r(v_\ell)]$$

- The task worth available to the players

$$w_r = \widetilde{w}_r \cdot (1 - q(r))$$



❖ Total Team Potential $\Phi(a)$

$$\Phi(a) = \sum_{r=1}^M \tilde{w}_r \cdot \left(1 - \prod_{v_\ell \in \{V\}_r} [1 - p_r(v_\ell)] \right)$$

where:

- $a = (a_{\mathcal{P}}, a_{\bar{\mathcal{P}}})$ is the joint action of both players $a_{\mathcal{P}}$ and non-players $a_{\bar{\mathcal{P}}}$
 - $\{V\}_r = \{\mathcal{P}\}_r \cup \{\bar{\mathcal{P}}\}_r$ is the set of all robots assigned to task r
- ❖ Once local players reach equilibrium $a_{\mathcal{P}}^*$, denote the new allocation for the whole team as $a^* = (a_{\mathcal{P}}^*, a_{\bar{\mathcal{P}}})$.

Remark: $\Phi(a)$ is the potential for the whole team; while $\phi(a_{\mathcal{P}})$ is the potential for the local players

Since the players and non-players are mixed and distributed over different tasks, how does the total team potential Φ change when local potential ϕ is increased?

Theorem^[1]: The optimal equilibrium a^* increases the total team potential $\Phi(a)$, i.e., $\Phi(a^*) \geq \Phi(a)$

Sketch of Proof:

$$\begin{aligned} \Phi(a) &= \sum_{r=1}^M \tilde{w}_r \cdot \left(1 - \prod_{\mathcal{P}_i \in \{\mathcal{P}\}_r} [1 - p_r(\mathcal{P})] \cdot \prod_{v_\ell \in \{\bar{\mathcal{P}}\}_r} [1 - p_r(v_\ell)] \right) \\ &= \sum_{r=1}^M \tilde{w}_r \cdot (1 - [1 - p(r)][1 - q(r)]) \\ &= \sum_{r=1}^M \tilde{w}_r [1 - q(r)] \cdot p(r) + \sum_{r=1}^M \tilde{w}_r \cdot q(r) \end{aligned}$$

❖ Joint success probability for players:
 $p(r) = 1 - \prod_{\mathcal{P}_i \in \{\mathcal{P}\}_r} [1 - p_r(\mathcal{P})]$

❖ Joint success probability for non-players:
 $q(r) = 1 - \prod_{v_\ell \in \{\bar{\mathcal{P}}\}_r} [1 - p_r(v_\ell)]$

Then, it can be written as

$$\begin{aligned} \Phi(a) &= \sum_{r=1}^M w_r \cdot p(r) + \sum_{r=1}^M \tilde{w}_r \cdot q(r) \\ &= \left(\sum_{r \in \tilde{\mathcal{A}}} w_r \cdot p(r) + \sum_{r \notin \tilde{\mathcal{A}}} w_r \cdot p(r) \right) + \sum_{r=1}^M \tilde{w}_r \cdot q(r) \\ &= \phi(a_{\mathcal{P}}) + \sum_{r \notin \tilde{\mathcal{A}}} w_r \cdot p(r) + \sum_{r=1}^M \tilde{w}_r \cdot q(r) \end{aligned}$$

Substitute: $w_r = \tilde{w}_r [1 - q(r)]$

Substitute: $\phi(a_{\mathcal{P}}) = \sum_{r \in \tilde{\mathcal{A}}} w_r \cdot p(r)$

The players should finish their current small left-over tasks

Not affected for non-players

The CARE Algorithm

The Reallocation Decision

Learning for equilibrium $a_{\mathcal{P}}^*$: Max-Logit can fast converge to the optimal equilibrium (i.e., maximizer to ϕ).

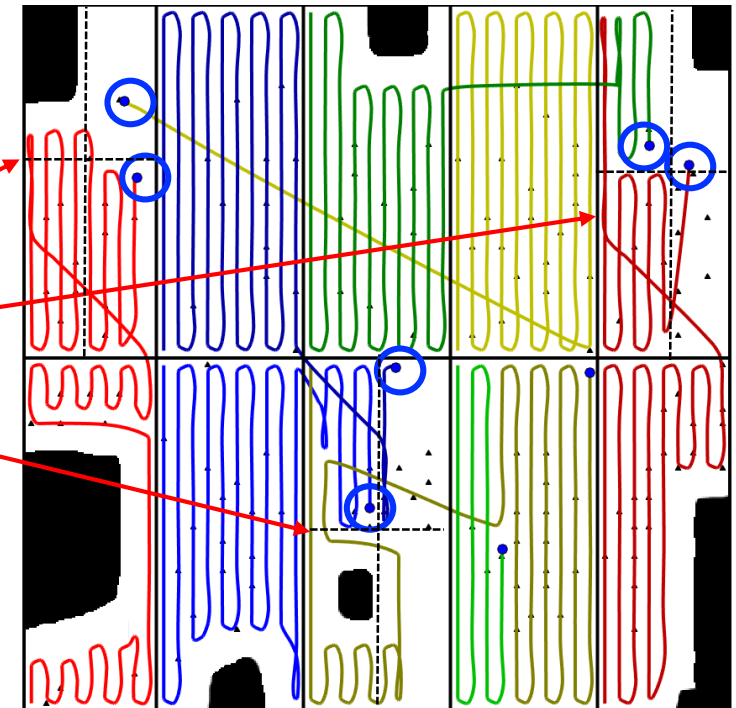
But... what if multiple robots are assigned to the same task, and/or what if there already are some non-player robots there?

Post-game Coordination

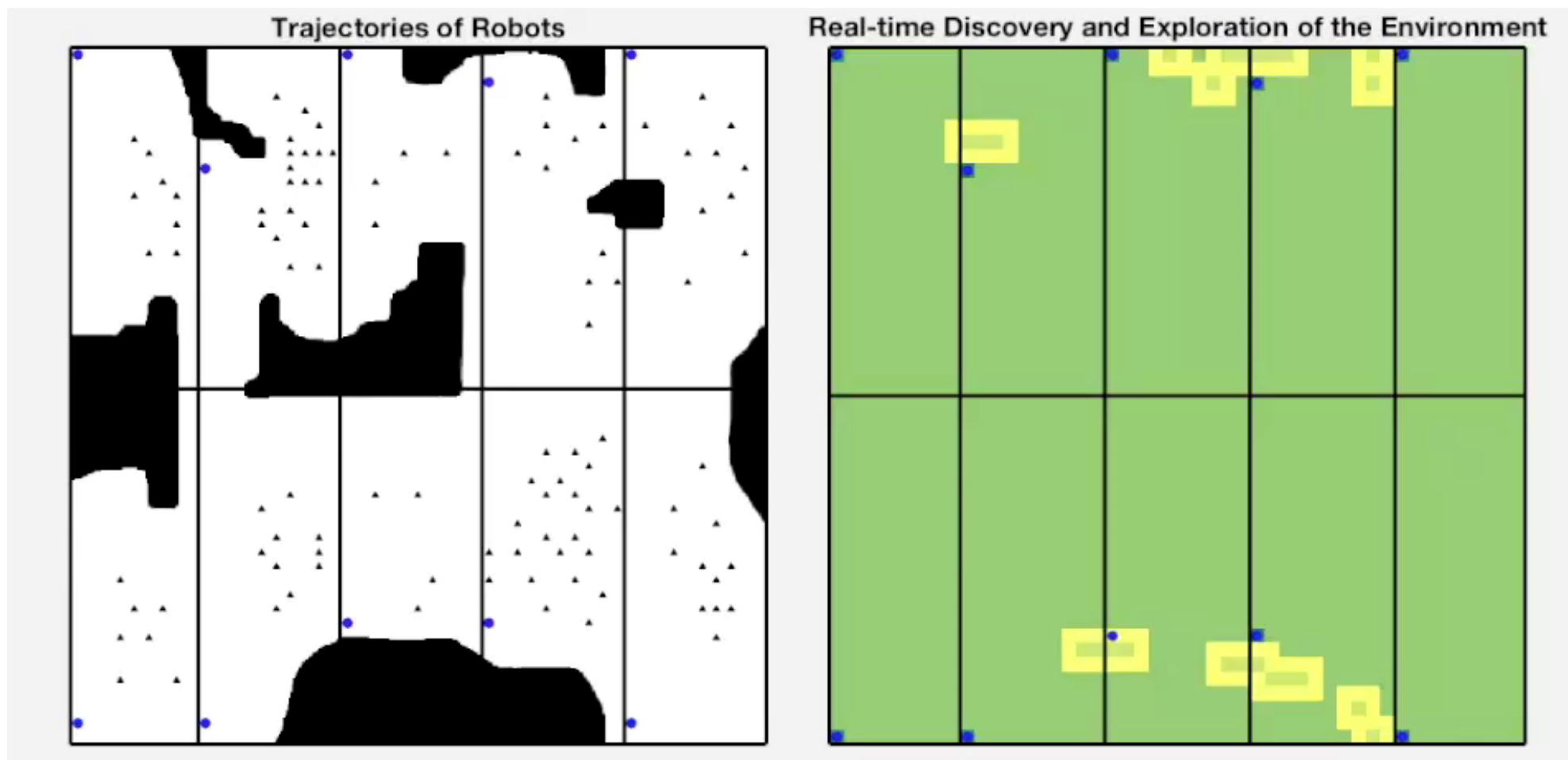
- First, evenly partition task r into $n_{\max} \in \mathbb{N}^+$ sub-regions
- Non-players continues search within the sub-region determined by its current location
- Each player $\mathcal{P}_i \in \{\mathcal{P}_i \in \mathcal{P} : a_i^* = r\}$ chooses the closest sub-region in an order based on its success probability $p_r(\mathcal{P}_i)$

Theorem 2^[1]: Complete coverage is guaranteed as long as at least one robot is still alive.

Division into $n_{\max} = 4$ sub-regions

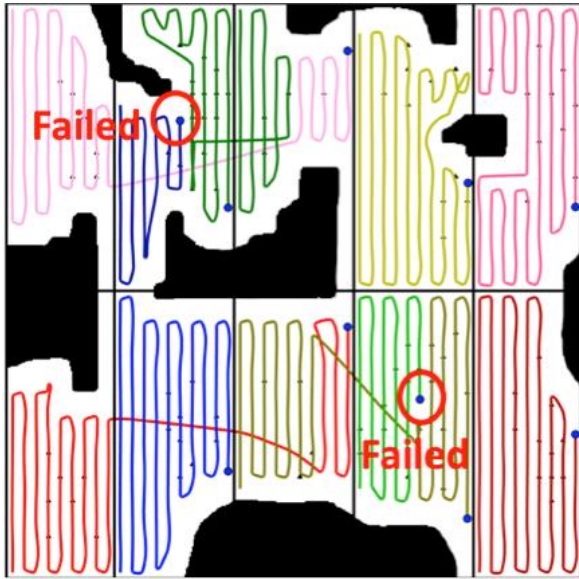


- ❖ The search area is of size $50\text{m} \times 50\text{m}$, partitioned into $M = 10$ tasks. A team of $N = 10$ robots are deployed.
- ❖ Vehicle: battery reliability parameters $\rho_0 \sim N(3 \times 10^{-3}, 7.5 \times 10^{-5})$ and $\rho_1 \sim N(1400, 35)$; laser range: 5m ; $u = 1\text{m/s}$; $\omega = 0.32$ cell/s
- ❖ Game parameters: $\kappa_1 = 6$, $\kappa_2 = 3$, $\eta = 30\text{s}$, $\gamma = 200\text{s}$, number of game computation cycles: 50

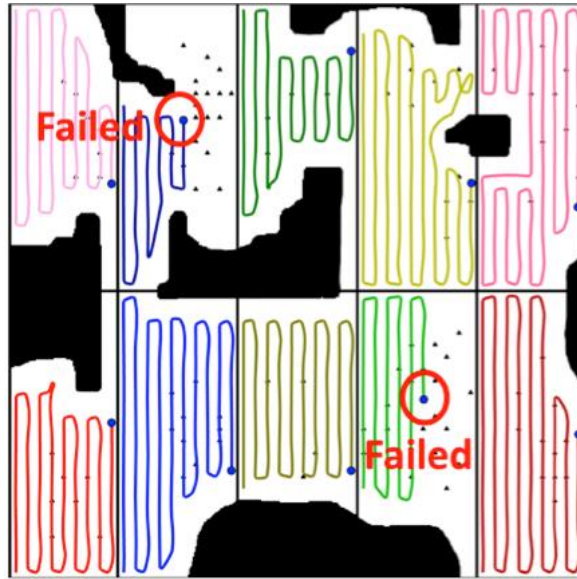


Coverage Trajectories using Alternative Methods

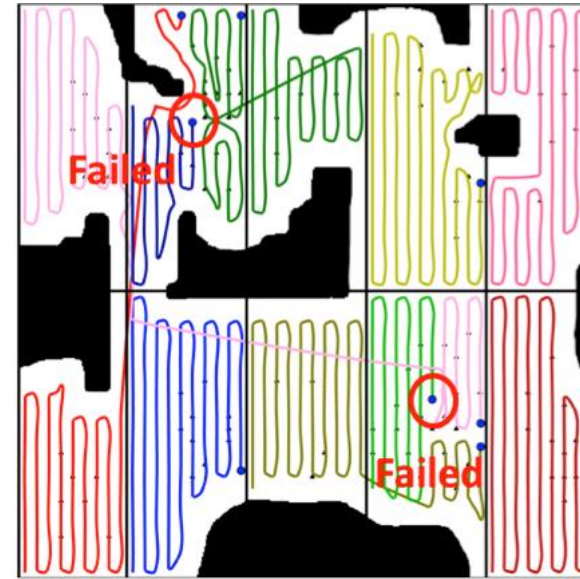
Incomplete coverage



(a) CARE

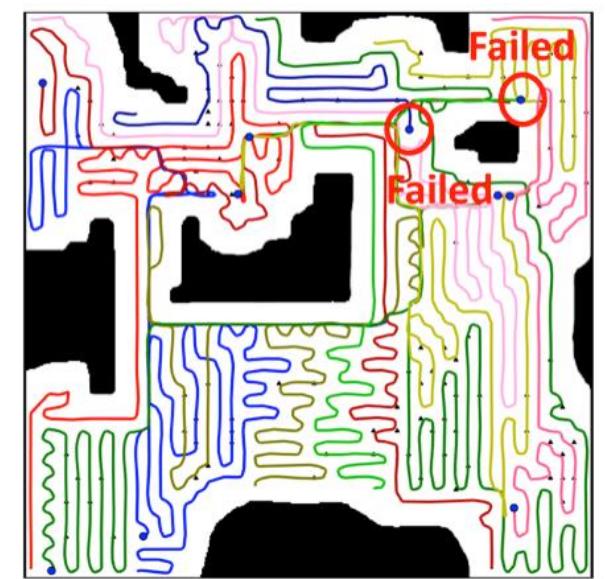


(b) Non-cooperative Coverage



(c) First-responder Coverage

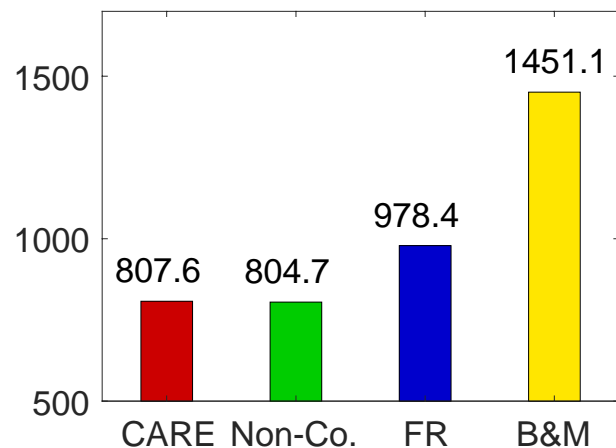
No task partition



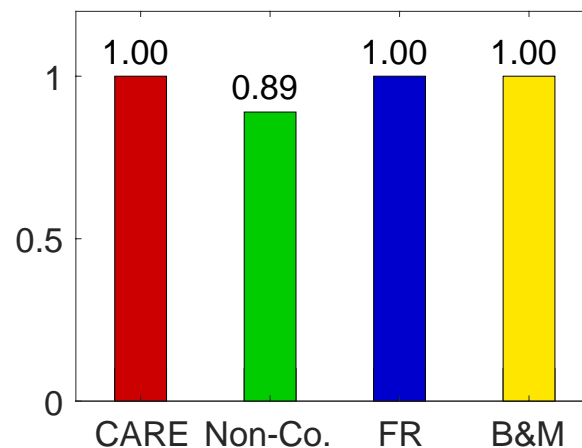
(c) Multi-robot Brick-and-Mortar

Performance Comparison with Alternative Methods

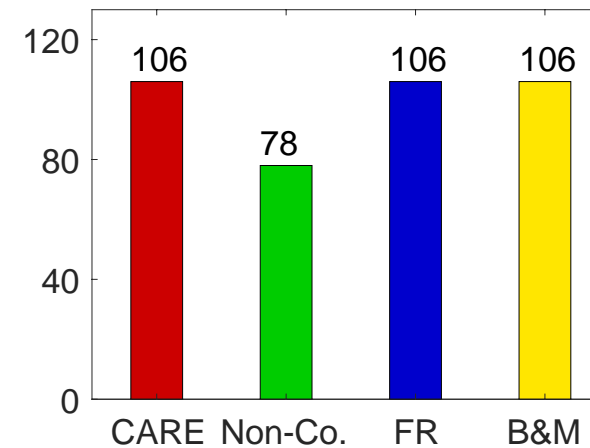
Coverage Time (CT)



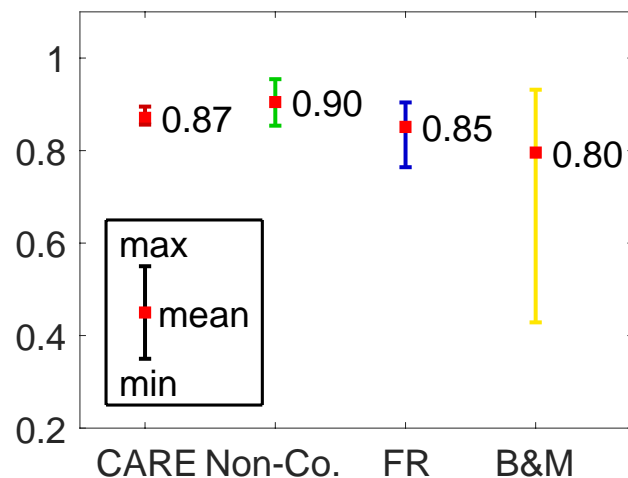
Coverage Ratio (CR)



Number of Targets Found ($NoTF$)



Remaining Reliability (RR)



Time of Target Discovery ($ToTD$)

